

Solutions - Midterm Exam

(February 20th @ 7:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (20 PTS)

- Compute the result of the following operations. The operands are signed fixed-point numbers. The result must be a signed fixed-point number. For the division, use $x = 5$ fractional bits.

1.0111 + 1.101001	1.010101 - 1000.0101	01.11111 + 0.10001
10.101 × 1.01101	1.001 × 0.1011	10.1010 ÷ 0.101

$$\begin{array}{r}
 \begin{array}{c} \text{C}_8=1 \\ \text{C}_7=1 \\ \text{C}_6=1 \\ \text{C}_5=1 \\ \text{C}_4=0 \\ \text{C}_3=0 \\ \text{C}_2=0 \\ \text{C}_1=0 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 1.0111 + \\ 1.101001 \\ \hline 1.000101 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{C}_{10}=1 \\ \text{C}_9=1 \\ \text{C}_8=1 \\ \text{C}_7=1 \\ \text{C}_6=1 \\ \text{C}_5=1 \\ \text{C}_4=1 \\ \text{C}_3=1 \\ \text{C}_2=0 \\ \text{C}_1=0 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 1.010101 - \\ 1000.0101 \\ \hline 0.011000 \end{array}
 \end{array}
 \quad
 \begin{array}{r}
 \begin{array}{c} \text{C}_8=0 \\ \text{C}_7=0 \\ \text{C}_6=1 \\ \text{C}_5=1 \\ \text{C}_4=1 \\ \text{C}_3=1 \\ \text{C}_2=1 \\ \text{C}_1=1 \\ \text{C}_0=0 \end{array} \\
 \begin{array}{r} 01.11111 + \\ 0.10001 \\ \hline 10.1010 \end{array}
 \end{array}$$

$$\begin{array}{r}
 10.101 \times \rightarrow 01.011 \times \rightarrow \begin{array}{r} 10011 \times \\ 1011 \\ \hline 10011 \\ 10011 \\ 00000 \\ 10011 \\ \hline 011010001 \\ \downarrow \\ 0.11010001 \end{array}
 \end{array}$$

$$\begin{array}{r}
 1.001 \times \rightarrow 0.111 \times \rightarrow \begin{array}{r} 1011 \times \\ 111 \\ \hline 1011 \\ 1011 \\ 1011 \\ \hline 01001101 \\ \downarrow \\ 0.1001101 \\ \downarrow \\ 1.0110011 \end{array}
 \end{array}$$

- ✓ $\frac{10.1010}{0.101}$: To unsigned (numerator) and then alignment, $a = 3$: $\frac{01.0110}{0.101} = \frac{01.011}{00.101} \equiv \frac{1011}{101}$

$$\begin{array}{r}
 001000110 \\
 101 \overline{) 101100000} \\
 \underline{101} \\
 01000 \\
 \underline{101} \\
 110 \\
 \underline{101} \\
 10
 \end{array}$$

Append $x = 5$ zeros: $\frac{101100000}{101}$

Integer Division:

$Q = 1000110, R = 10$

$\rightarrow Qf = 10.00110 (x = 5)$

Final result (2C): $\frac{10.1010}{0.101} = 2C(010.0011) = 101.1101$

PROBLEM 2 (10 PTS)

- Represent these numbers in Fixed Point Arithmetic (signed numbers). Select the minimum number of bits in each case.

✓ -32.125

✓ 17.75

$+32.125 = 0100000.001 \Rightarrow -32.125 = 1011111.111$

$+17.25 = 010001.11$

- Complete the table for the following fixed point formats (signed numbers): (6 pts.)

Integer bits	Fractional Bits	FX Format	Range	Resolution
8	6	[14 6]	$[-2^7, 2^7 - 2^{-6}]$	2^{-6}
6	4	[10 4]	$[-2^5, 2^5 - 2^{-4}]$	2^{-4}

PROBLEM 3 (40 PTS)

- Calculate the result (provide the 32-bit result) of the following operations with 32-bit floating point numbers. Truncate the results when required. When doing fixed-point division, use 4 fractional bits. Show your procedure.

✓ 40D00000 + C2EA0000	✓ 50A90000 - 4F480000	✓ 80200000 × 7AB80000	✓ FB380000 ÷ 48C00000
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- ✓ $X = 40D00000 + C2EA0000$:

40D00000: 0100 0000 1101 0000 1000 0000 0000 0000

$$e + \text{bias} = 10000001 = 129 \rightarrow e = 129 - 127 = 2$$

$$\text{Significand} = 1.101$$

$$40D00000 = 1.101 \times 2^2$$

C2EA0000: 1100 0010 1110 1010 0000 0000 0000 0000

$$e + \text{bias} = 10000101 = 133 \rightarrow e = 133 - 127 = 6$$

$$\text{Significand} = 1.110101$$

$$C2EA0000 = -1.110101 \times 2^6$$

$$X = 1.101 \times 2^2 - 1.110101 \times 2^6 = +\frac{1.101}{2^4} \times 2^6 - 1.110101 \times 2^6 = (0.0001101 - 1.110101) \times 2^6$$

To subtract these numbers, we first convert to 2C:

$$R = 0.0001101 - 01.110101 = 0.0001101 + 10.001011 \text{ (2C addition)}$$

$$\text{The result in 2C is: } R = 10.0100011, -R = 01.1011101$$

$$* \text{ Note that you can also do unsigned subtraction: } X = -(1.110101 - 0.0001101 -) \times 2^6$$

For floating point, we need to convert to sign-and-magnitude:

$$\Rightarrow R(SM) = -1.1011101$$

$$\begin{array}{r} \overset{C_9}{0} \overset{C_8}{0} \overset{C_7}{0} \overset{C_6}{0} \overset{C_5}{1} \overset{C_4}{1} \overset{C_3}{1} \overset{C_2}{0} \overset{C_1}{0} \overset{C_0}{0} \\ 0 \ 0.0 \ 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ + \\ 1 \ 0.0 \ 0 \ 1 \ 0 \ 1 \ 1 \ 0 \\ \hline 1 \ 0.0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 1 \end{array}$$

$$X = -1.1011101 \times 2^6, e + \text{bias} = 6 + 127 = 133 = 10000101$$

$$X = 1100 \ 0010 \ 1101 \ 1101 \ 0000 \ 0000 \ 0000 \ 0000 = C2DD0000$$

- ✓ $X = 50A90000 - 4F480000$:

50A90000: 0101 0000 1010 1001 0000 0000 0000 0000

$$e + \text{bias} = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

$$\text{Significand} = 1.0101001$$

$$50A90000 = 1.0101001 \times 2^{34}$$

4F480000: 0100 1111 0100 1000 0000 0000 0000 0000

$$e + \text{bias} = 10011110 = 158 \rightarrow e = 158 - 127 = 31$$

$$\text{Significand} = 1.1001$$

$$4F480000 = 1.1001 \times 2^{31}$$

$$X = 1.0101001 \times 2^{34} - 1.1001 \times 2^{31} = 1.0101001 \times 2^{34} - \frac{1.1001}{2^3} \times 2^{34}$$

$$X = (1.0101001 - 0.0011001) \times 2^{34} \text{ (unsigned subtraction)}$$

$$\begin{array}{r} \overset{b_8}{0} \overset{b_7}{0} \overset{b_6}{0} \overset{b_5}{1} \overset{b_4}{0} \overset{b_3}{0} \overset{b_2}{0} \overset{b_1}{0} \overset{b_0}{0} \\ 1.0 \ 1 \ 0 \ 1 \ 0 \ 0 \ 1 \ - \\ 0.0 \ 0 \ 1 \ 1 \ 0 \ 0 \ 1 \\ \hline 1.0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0 \end{array}$$

$$X = 1.001 \times 2^{34}, e + \text{bias} = 34 + 127 = 161 = 10100001$$

$$X = 0101 \ 0000 \ 1001 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = 50900000$$

- ✓ $X = 80200000 \times 7AB80000$:

80200000: 1000 0000 0010 0000 0000 0000 0000 0000

$$e + \text{bias} = 00000000 = 0 \rightarrow \text{Denormal number} \rightarrow e = -126 \quad \text{Significand} = 0.01$$

$$80200000 = -0.01 \times 2^{-126}$$

7AB80000: 0111 1010 1011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110101 = 245 \rightarrow e = 245 - 127 = 118$$

$$\text{Significand} = 1.0111$$

$$7AB80000 = 1.0111 \times 2^{118}$$

$$X = (-0.01 \times 2^{-126}) \times (1.0111 \times 2^{118}) = -0.010111 \times 2^{-8} = -1.0111 \times 2^{-10}$$

$$e + \text{bias} = -10 + 127 = 117 = 01110101$$

$$X = 1011 \ 1010 \ 1011 \ 1000 \ 0000 \ 0000 \ 0000 \ 0000 = BAB80000$$

- ✓ $X = FB380000 \div 48C00000$:

FB380000: 1111 1011 0011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110110 = 246 \rightarrow e = 246 - 127 = 119$$

$$\text{Significand} = 1.0111$$

$$FB380000 = -1.0111 \times 2^{119}$$

48C00000: 0100 1000 1100 0000 0000 0000 0000 0000

$$e + \text{bias} = 10010001 = 145 \rightarrow e = 145 - 127 = 18$$

$$\text{Significand} = 1.1$$

$$48C00000 = 1.1 \times 2^{18}$$

$$X = -\frac{1.0111 \times 2^{119}}{1.1 \times 2^{18}} = -\frac{1.0111}{1.1} \times 2^{101}$$

Alignment:

$$\frac{1.0111}{1.1} = \frac{1.0111}{1.1000} = \frac{10111}{11000}$$

Append $x = 4$ zeros: $\frac{101110000}{11000}$

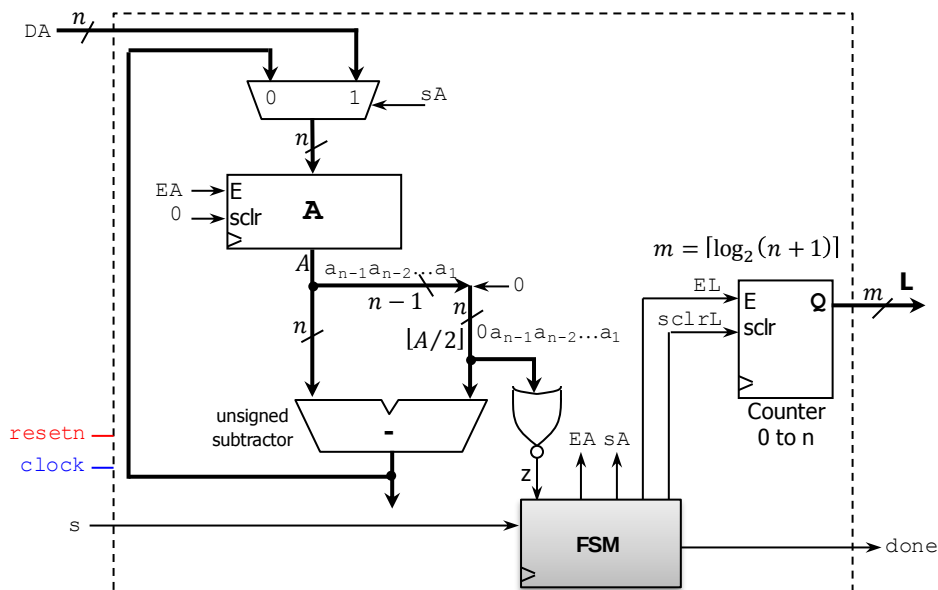
Integer division
 $Q = 1111 \rightarrow Qf = 0.1111$

Thus: $X = -0.1111 \times 2^{101} = -1.111 \times 2^{100}$
 $e + \text{bias} = 100 + 127 = 227 = 11100011$

$X = 1111 \ 0001 \ 1111 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = F1F00000$

PROBLEM 4 (30 PTS)

- "Integer Base-2 Logarithm": This circuit computes $L = \lfloor \log_2 A \rfloor$, where A is an unsigned n -bit number stored in register A . The digital system is depicted below: FSM + Datapath. Example: For $n = 8$: if $A = 00110110$, then $L = 0110$.
 - ✓ Note that for $n = 8$, if $A = a_7a_6a_5a_4a_3a_2a_1a_0$, then $\lfloor A/2 \rfloor = 0a_7a_6a_5a_4a_3a_2a_1$. Also: $z = 1$ if $\lfloor A/2 \rfloor = 0$.
 - ✓ m-bit counter: $sclr$. If $E = sclr = 1$, the count is initialized to zero. If $E = 1, sclr = 0$, the count is increased by 1.
 - ✓ Register: If $E = 1: sclr = 1 \rightarrow \text{Clear}, sclr = 0 \rightarrow \text{Load}$.
- Sketch the Finite State Machine diagram (in ASM form) given the algorithm (for $n = 8, m = 4$). (18 pts.)
 - ✓ The process begins when s is asserted, at this moment we capture DA on register A . Then the process continues by updating A and it is concluded when $A = 1$. The signal $done$ is asserted when we finish computing $\lfloor \log_2 A \rfloor$.
- Complete the timing diagram where $n = 8, m = 4$. (12 pts.)



ALGORITHM

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L ← 0
while A ≠ 1
    A ← A - ⌊A/2⌋
    L ← L + 1
end while
  
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Finite State Machine:

